



- Project 4 Simulation of Continuous Dynamic Partition Allocation
- Assume that in the initial state, the available memory space is 640KB, and there is the following request sequence:
- •job1 applies for 130KB
- •job2 applies for 60KB
- •job3 applies for 100KB
- •job2 releases 60KB
- •job4 applies for 200KB

- •job3 releases 100KB
- •job1 releases 130KB
- •job5 applies for 140KB
- •job6 applies for 60KB
- •job7 applies for 50KB
- •job6 releases 60KB



Task

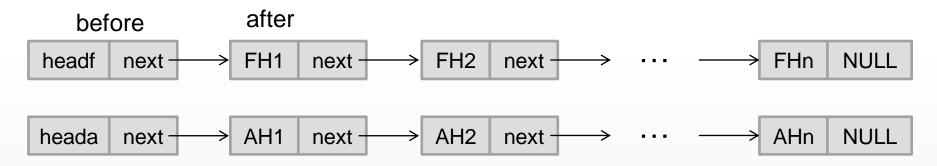
- The allocation process alloc() and the recycling process free() for dynamic partition are respectively implemented based on the first-fit algorithm and the best-fit algorithm in C language.
- Free holes are managed through a free hole chain.
- When allocating memory, the system gives priority to using the lower space of the free area.
- Shows the status of free hole chains after each allocation and recycling.



- Free and allocated hole chain
- struct node // node for the chain
- •
- int address, size;
- struct node *next;
- **■** };
- typedef struct node RECT;
- Example



- Free and allocated hole chain
- Free hole(FH)
- Allocated hole(AH)





Pseudocode for traversing free hole chain for first-fit before=headf; after=headf->next; while(after!=NULL){ if(after->size > the requested size){ allocate; break; }else if(after->size == the requested size){ allocate; before-next=after->next; free(after); }else{ before=before->next; after=after->next;



- Pseudocode for traversing free hole chain to reclaim memory
- before=headf;
- after=headf->next;
- back=malloc(sizeof(RECT));
- back->address=the reclaimed address;
- back->size=the reclaimed size;
- flag=false;



```
while(!flag){
- If(after==NULL){
- before->next=back;
- back-next=after;
- flag=true;
- break;
- }
```



```
if(before->address+before->size==back->address){
  if(back->address+back->size==after->address){// merge back with before and after
   if(before!=headf){
     before->size=before->size+back->size+after->size;
     before->next=after->next;
     free(after);
     free(back);
     flag=true;
     break;
  }else{ //merge back with before
    if(before!=headf){
      before->size=before->size+back->size;
     free(back);
     flag=true;
     break;
```



```
if(back->address+back->size==after->address){// merge back with after
     after->size=after->size+back->size:
     after->address=back->address;
     free(back);
     flag=true;
     break:
if((before->address<=back->address)&&(after->address>back->address)){
  // insert back to the free hole chain
 before->next=back;
 back->next=after;
 flag=true;
}else{
 before=before->next;
 after=after->next;
}//while
```



- Pseudocode for traversing allocated hole chain to reclaim memory
- before=heada;
- after=heada->next;
- while(after!=NULL){

```
If(after->size == the reclaimed size && after->address==the reclaimed address){
    before->next=after->next;
    free(after);
    break;
    }else{
        before=before->next;
        after=after->next;
    }
}
```





- Another request sequence:
 - •job1 applies for 130KB
 - •job2 applies for 60KB
 - •job3 applies for 100KB
 - •job1 releases 130KB
 - •job4 applies for 200KB
 - •job3 releases 100KB
 - •job5 applies for 150KB
 - •job1 releases 60KB
 - •job6 applies for 290KB
 - •job5 releases 150KB
 - •job6 releases 290KB
 - •job4 releases 200KB